

(2.55) Remark.

Let's analyze the cost of solving the system $Ax = b$, with $A \in \mathbb{R}^{n \times n}$ an invertible matrix, using the EGPP procedure. The procedures executed are EGPP, SA, and SI. We have:

$$C(\text{EGPP}) = \frac{2}{3} n^3 + \dots, \quad C(\text{SA}) = C(\text{SI}) = n^2$$

Note that while in the computation of SA and SI *only* arithmetic operations are performed, in the calculation of EGPP *comparisons* are also performed, but their number is *negligible* compared to that of the arithmetic operations.¹

(Exercise: Determine the number of comparisons performed by EGPP.)

The overall number of arithmetic operation is:

$$C(\text{EGPP}) + C(\text{SA}) + C(\text{SI}) = \frac{2}{3} n^3 + \dots$$

When we use the qr procedure to solve the system, the qr, matrix by column product (denoted by: pmc) and SI procedures are performed. We have:

$$C(\text{qr}) = \frac{4}{3} n^3 + \dots, \quad C(\text{pmc}) = 2 n^2 + \dots, \quad C(\text{SI}) = n^2$$

Note that in the qr procedure (as in the GS procedure) the computation of square roots is also performed but their number (n) is *negligible* compared to that of the arithmetic operations.

(Exercise: Determine the number of arithmetic operations performed by pmc.)

The overall number of arithmetic operation is:

$$C(\text{qr}) + C(\text{pmc}) + C(\text{SI}) = \frac{4}{3} n^3 + \dots$$

The dominant term in the arithmetic cost of the procedure using qr is therefore *two times larger* than that of the procedure using EGPP.

¹ It can be reasonably assumed that the time needed to compare two machine numbers is *similar* to the time needed to perform an arithmetic operation on the same numbers.

(2.4) ITERATIVE METHODS FOR THE SOLUTION OF A SYSTEM OF LINEAR EQUATIONS

(2.56) Definition (iterative method for the solution of a system of linear equations).

Let $H \in \mathbb{R}^{n \times n}$ and $c \in \mathbb{R}^n$. The iterative method defined by H and c is the application that maps each vector $g \in \mathbb{R}^n$ to the sequence of vectors $x(k)$ defined by:

$$x(0) = g, \quad x(k) = H x(k-1) + c \quad \text{for } k = 1, 2, \dots$$

(2.57) Remark.

- The iterative method defined by H and c is the iterative method defined by the function $h: \mathbb{R}^n \rightarrow \mathbb{R}^n$ such that:

$$h(x) = Hx + c$$

The function h is *continuous* therefore (see Remark (1.54) in Lecture 8) if the sequence $x(k)$ generated by the method is convergent, then its limit $v \in \mathbb{R}^n$ is a *fixed point* of h , that is, it satisfies the relation:

$$v = H v + c \quad \text{equivalent to} \quad (I - H) v = c$$

and this last relationship means that:

v is the solution of the system of linear equations $(I - H) x = c$

- Let $A \in \mathbb{R}^{n \times n}$ be an invertible matrix. The iterative method defined by H and c *can be used* to approximate the solution of the system $A x = b$ if:
 - (1) the systems $A x = b$ and $(I - H) x = c$ are *equivalent* (in particular: $I - H$ is invertible) and
 - (2) it is (practically) possible to find $g \in \mathbb{R}^n$ from which the sequence generated by the method is convergent.

(2.58) Example.

(1) Let:

$$A = \begin{bmatrix} 1/2 & 0 \\ 0 & -1 \end{bmatrix}, \quad b = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

- Set: $H = I - A = \begin{bmatrix} 1/2 & 0 \\ 0 & 2 \end{bmatrix}$ and $c = b$; the systems $A x = b$ e $(I - H) x = c$ are equivalent;
- Let $g = \begin{bmatrix} g_1 \\ g_2 \end{bmatrix}$. The sequence generated by the method defined by H and c is then:

$$x(0) = g, \quad x(1) = H x(0) + c = H g, \quad x(2) = H x(1) + c = H^2 g, \quad \dots$$

hence:

$$x(k) = H^k g = \begin{bmatrix} (1/2)^k & 0 \\ 0 & 2^k \end{bmatrix} g = \begin{bmatrix} (1/2)^k g_1 \\ 2^k g_2 \end{bmatrix}$$

The sequence is convergent *if and only if* $g_2 = 0$. In this case we have:

$$\lim_{k \rightarrow \infty} x(k) = 0$$

and the sequence converges to the unique solution of the system $A x = b$.

(2) Let:

$$A = \begin{bmatrix} 2 & 1 \\ 1 & 2 \end{bmatrix}, \quad b = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

- Set: $J = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$ and rewrite $A = 2 I + J$. Then:

$$A x = b \quad \text{is equivalent to} \quad x = - (1/2) J x + (1/2) b$$

i.e., defined $H = -(1/2) J$ and $c = (1/2) b$:

$$A x = b \quad \text{is equivalent to} \quad (I - H) x = c$$

- The eigenvalues of the matrix H are: $\lambda_1 = -1/2$ and $\lambda_2 = 1/2$, hence H is a diagonalizable matrix.² It is:

$$H = S \begin{bmatrix} -1/2 & 0 \\ 0 & 1/2 \end{bmatrix} S^{-1} \quad \text{where} \quad S = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

- Set $g = \begin{bmatrix} g_1 \\ g_2 \end{bmatrix}$. The sequence generated by the method defined by H and c is:

$$x(k) = H^k g = S \begin{bmatrix} (-1/2)^k & 0 \\ 0 & (1/2)^k \end{bmatrix} S^{-1} g$$

i.e., defined $y = S^{-1} g$:

$$x(k) = \begin{bmatrix} (-1/2)^k y_1 \\ (1/2)^k y_2 \end{bmatrix}$$

In this case we have:

$$\text{for every } g \in \mathbb{R}^2 : \lim_{k \rightarrow \infty} x(k) = 0$$

that is: for every $g \in \mathbb{R}^2$ the sequence converges to the unique solution of the system $A x = b$.

2 Recall that (1) a matrix $M \in \mathbb{R}^{n \times n}$ is diagonalizable if there exist a diagonal matrix Λ and an invertible matrix S such that: $MS = S\Lambda$, or $M = S\Lambda S^{-1}$; the elements $\lambda_1, \dots, \lambda_n$ on the diagonal of Λ are the eigenvalues of M , the k -th column of S is an eigenvector associated with the eigenvalue λ_k ; (2) if a matrix M has distinct eigenvalues then M is diagonalizable.

(3) Let $A = -I$ and $b = 0$.

- Set $H = I - A$ and $c = b$; the systems $Ax = b$ e $(I - H)x = c$ are equivalent.
- The sequence generated by the iterative method defined by H and c starting from $g \in \mathbb{R}^n$ is:

$$x(k) = H^k g = 2^k g$$

The sequence is convergent *if and only if* $g = 0$ and, in this case, it converges to the unique solution of the system $Ax = b$.